

Product Programs in the Wild: Retrofitting Program Verifiers to Check Information Flow Security

Marco Eilers
Severin Meier
Peter Müller

Noninterference

```
def compute(l: Int, h: Int) -> Int:  
  var res;  
  if (h > 0):  
    res := 1  
  else:  
    res := 2  
  return res;  
}
```

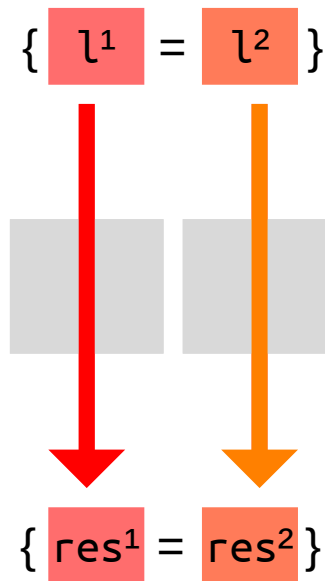
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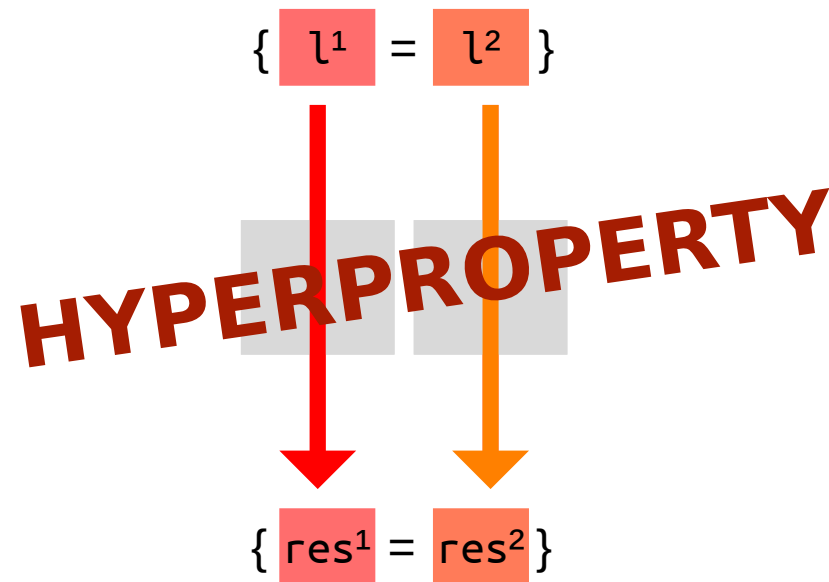
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Self-composition and Product Programs

[Barthe et al., Math. Struc. Comp. Sci. 21(6), 2011]

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Self-composition and Product Programs

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```
var res1;  
res1 := foo(l1)  
...
```

```
var res2;  
res2 := foo(l2)  
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Self-composition and Product Programs

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{ l1 == l2 }  
  
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var res2;  
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Modular Product Programs

[Eilers, Müller, Hitz, ESOP 2018]

```
{ l1 == l2 }  
var res1;  
var res2;  
x1,x2 := foo(y1, y2)  
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Modular Product Programs

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- Calls *not* duplicated

Modular Product Programs

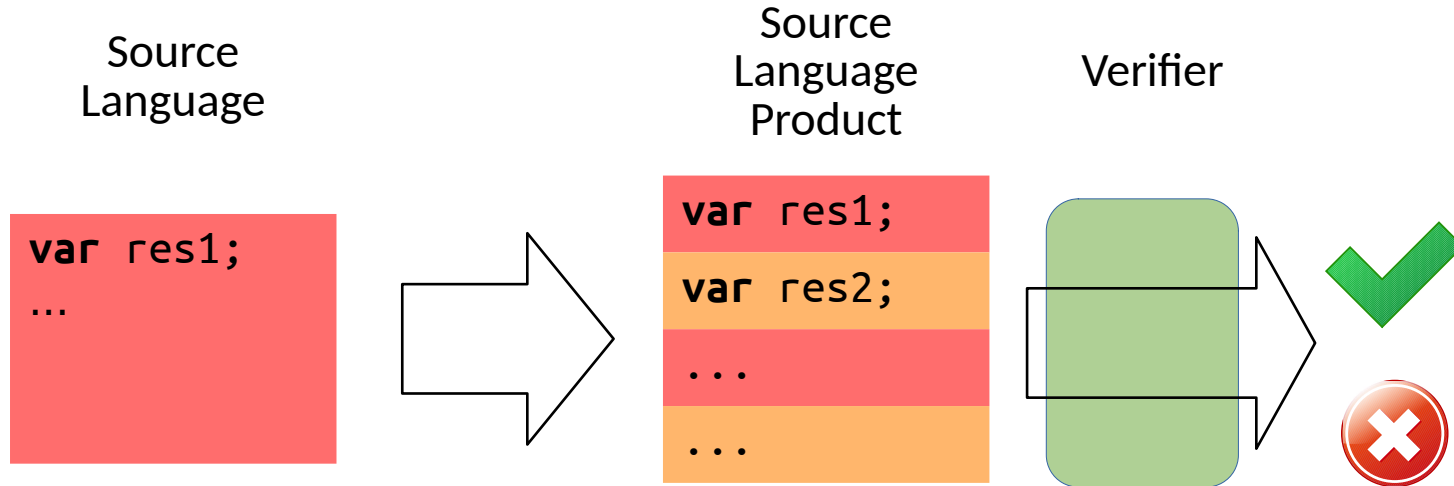
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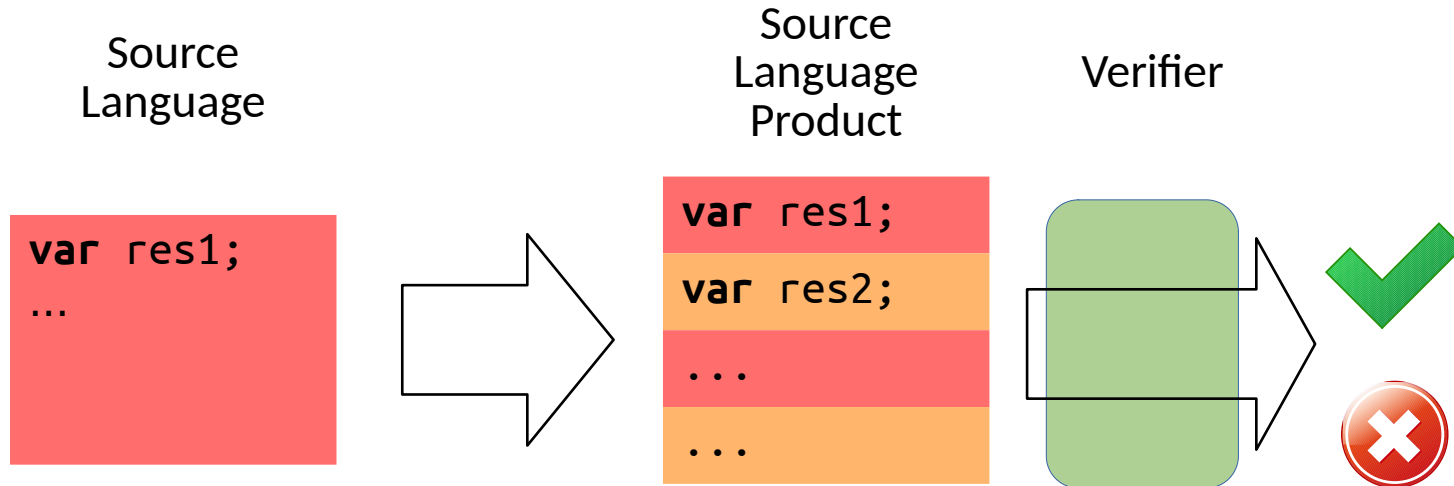
```
def foo(y1, y2)  
  # ensures:  
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  ...
```

- Calls *not* duplicated
- Enables use of *relational* specifications

Problem: Real Languages

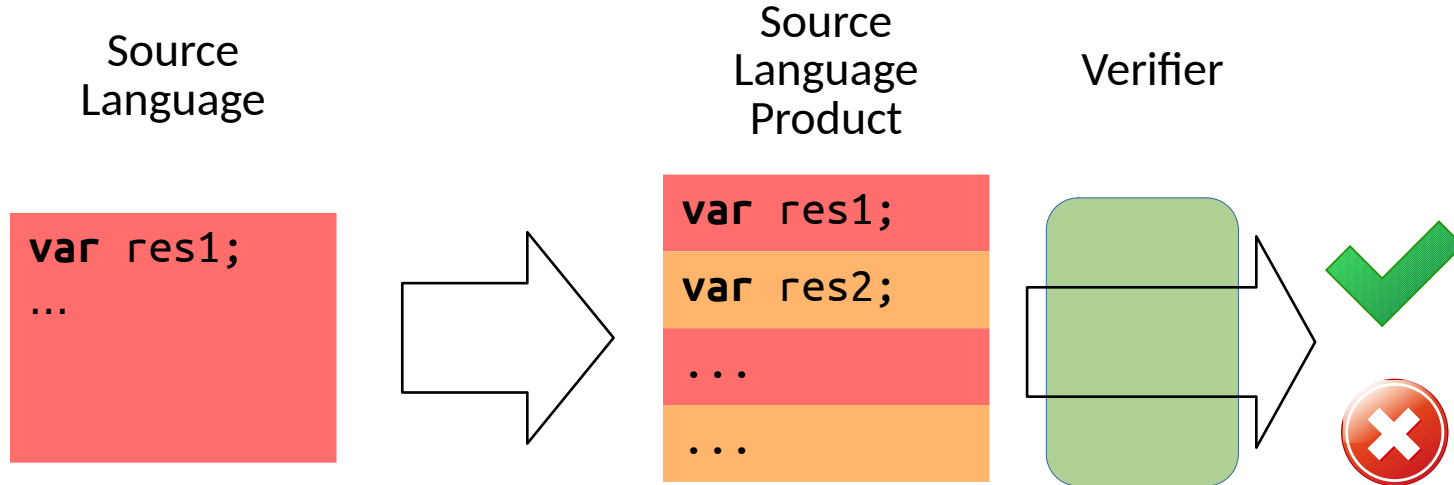


Problem: Real Languages



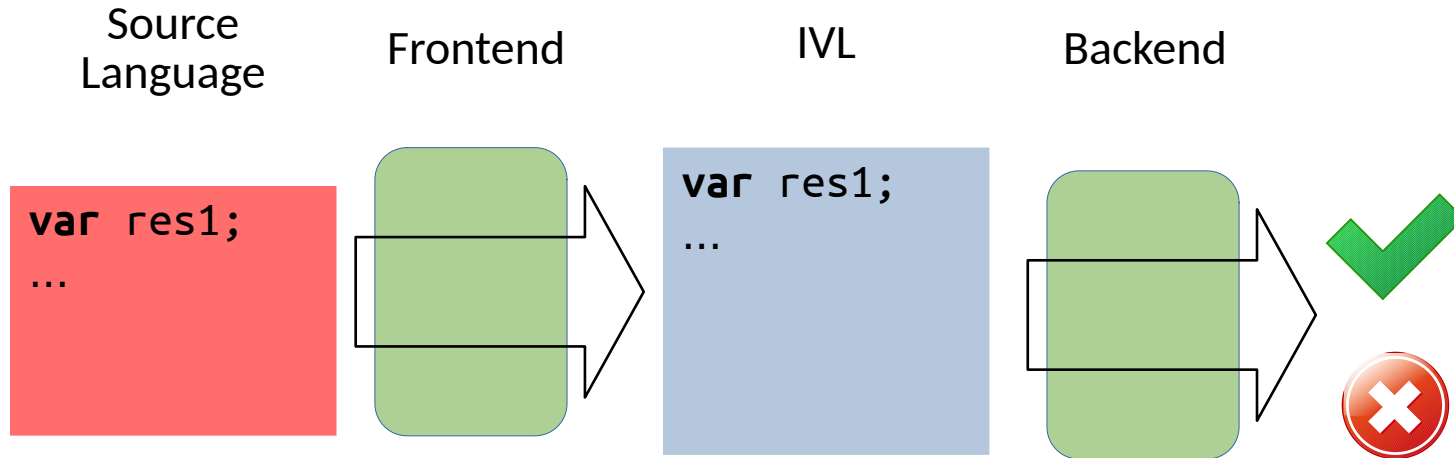
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- Specific to source language

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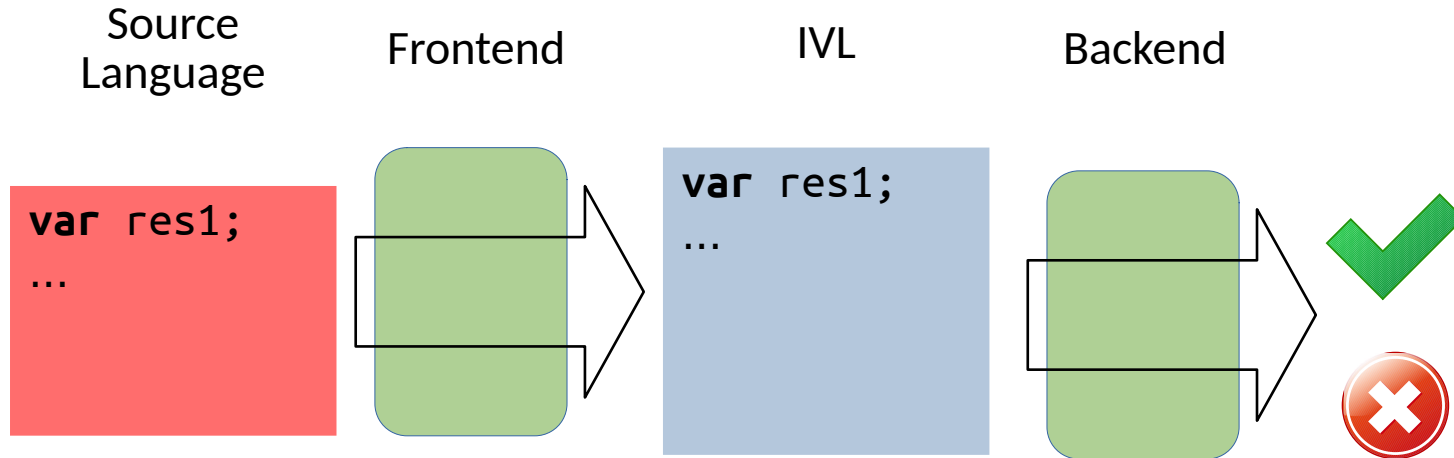


- Complex definition
- Specific to source language
- Only sequential

Typical Verifier Architecture

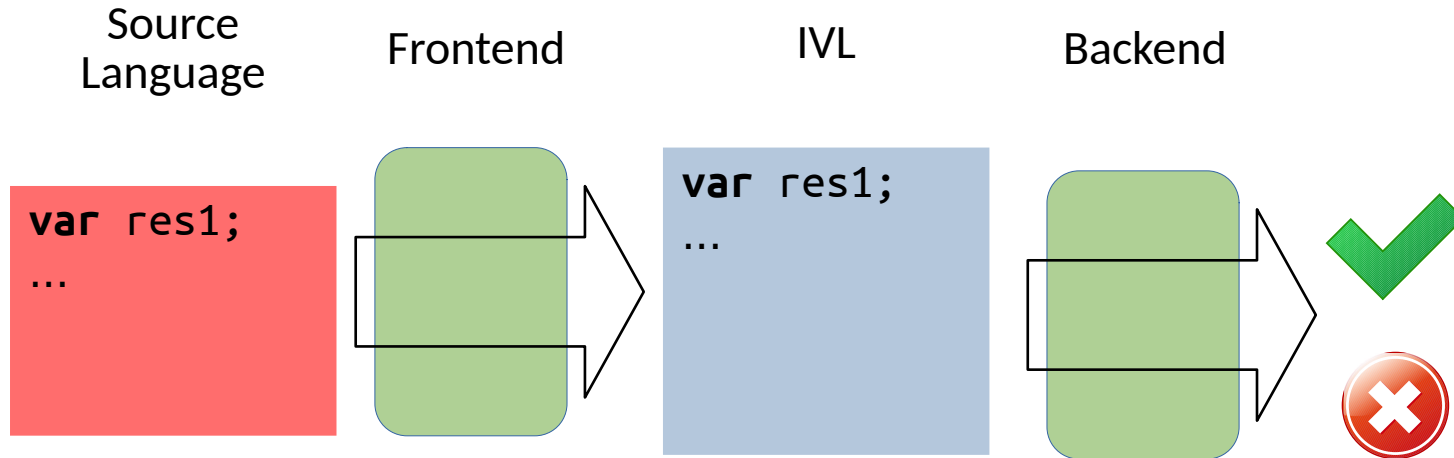


Typical Verifier Architecture



- Simple, sequential
- Examples: Boogie, Viper, Why3

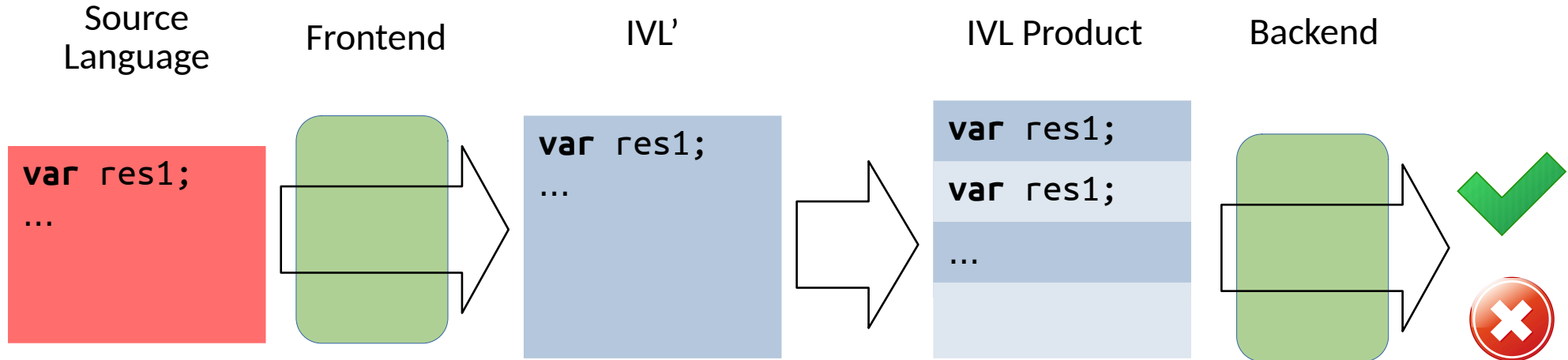
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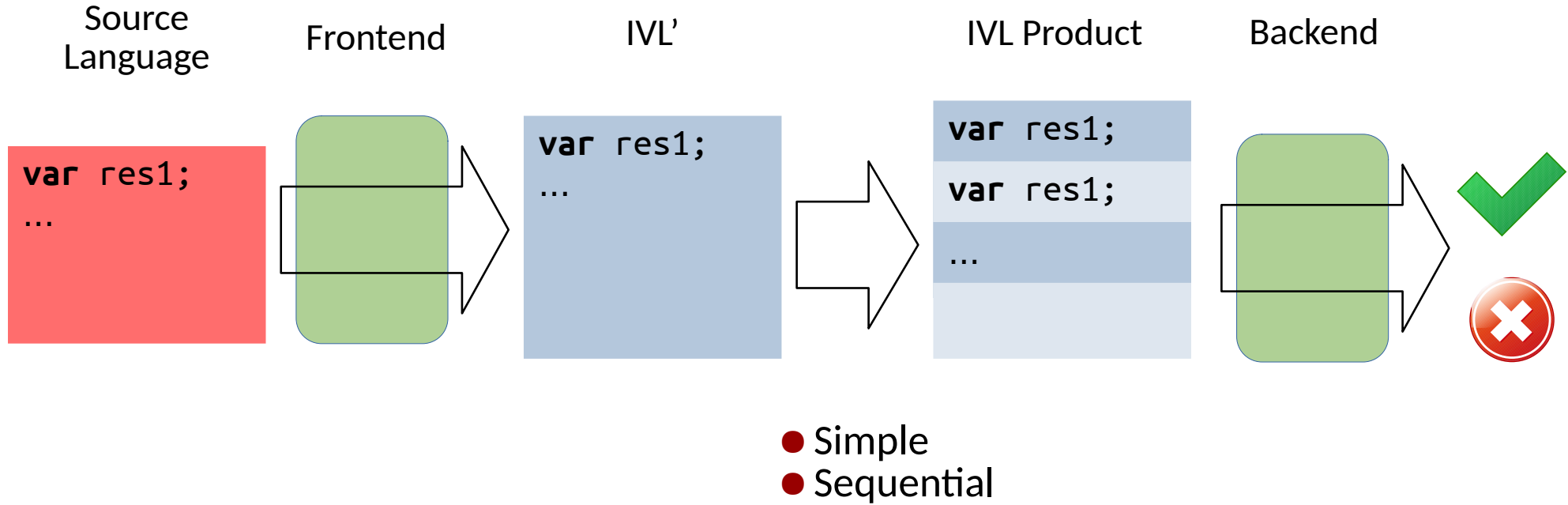
- Boogie: Dafny, VCC, GPUVerify
- Viper: Vercors, Prusti, Nagini
- Why3: Frama-C

- Simple, sequential
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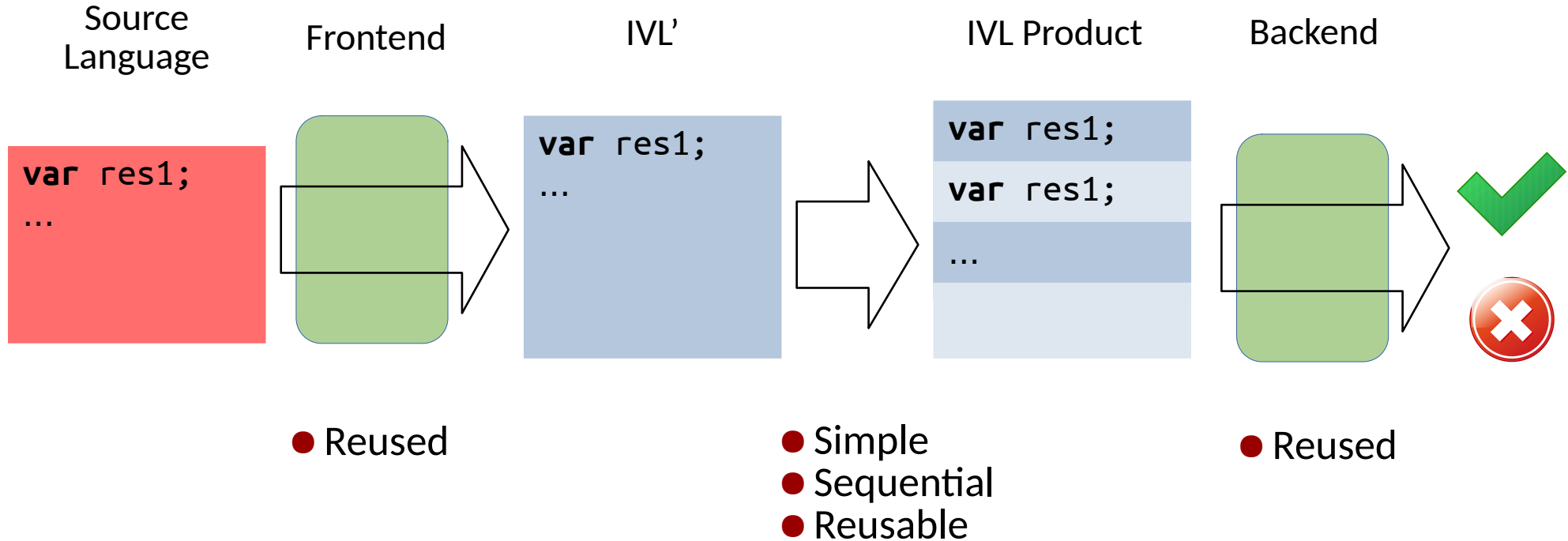
Proposed Architecture



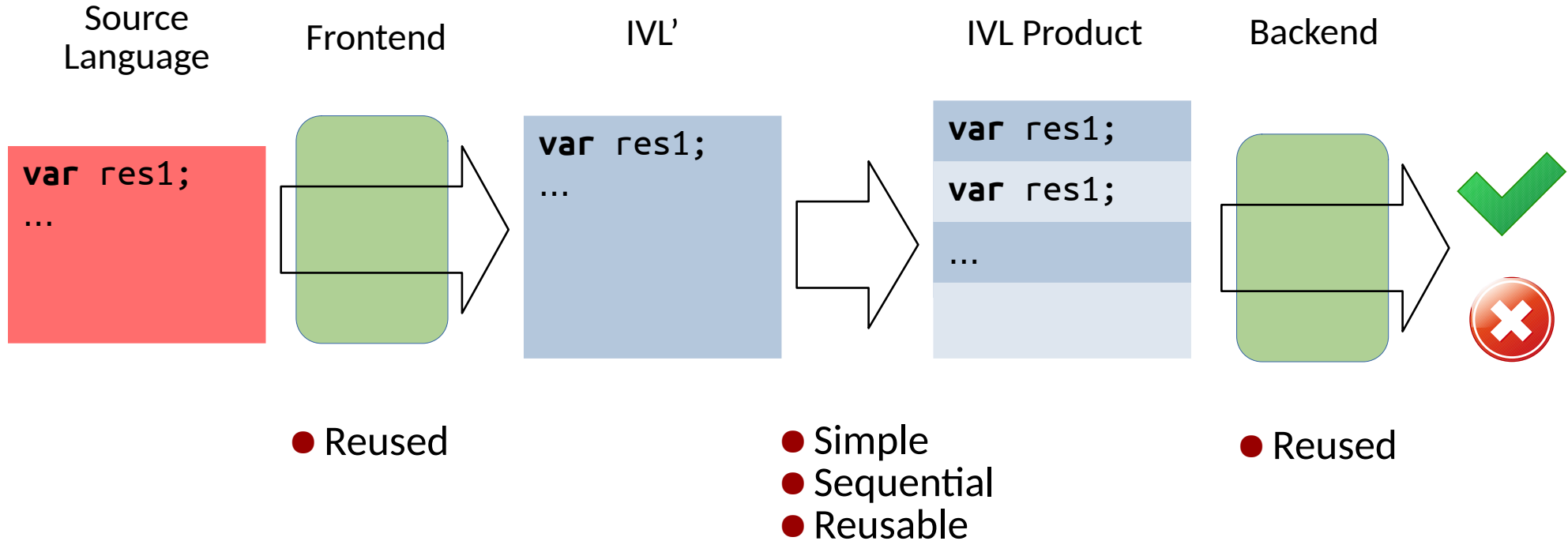
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Proposed Architecture



Already used by SymDiff [Lahiri et al., CAV 2012]

Questions

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What are the requirements to ensure soundness?



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How can this approach be extended to concurrent programs?

$\vdash$
 $c \parallel c$

Questions



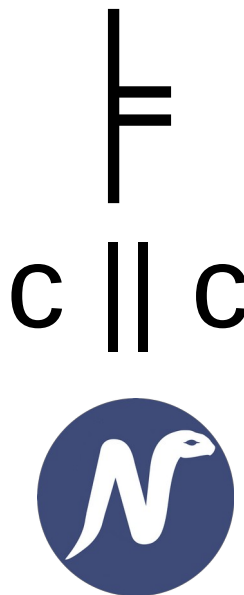
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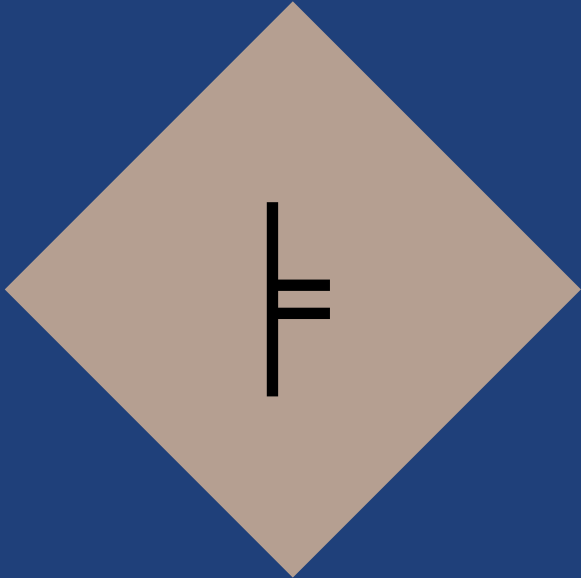
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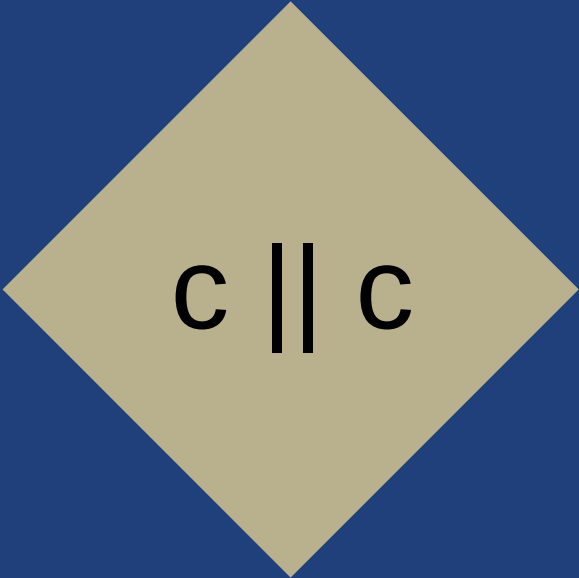
Does this approach allow for expressive and performant verification in practice?



Overview



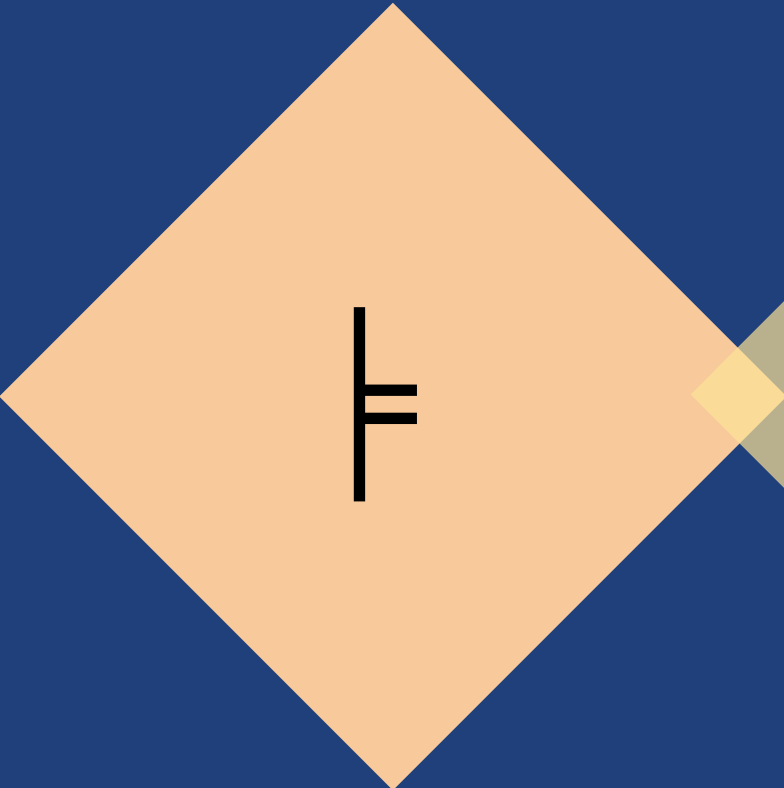
$\models$



$c \parallel c$



Soundness

A large orange diamond shape, tilted 45 degrees, with a smaller yellow diamond overlapping its right side.

$\models$

A small yellow diamond shape, tilted 45 degrees, overlapping the right side of the orange diamond and the left side of the olive diamond.

$c \parallel c$



Soundness Problem (1): Typical Encoding

```
def bar(a: A):  
    # ensures result != 0  
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```
class A:  
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        return 1  
  
class B(A):  
    def foo(self):  
        # ensures result >= 2  
        return 2
```

Soundness Problem (1): Typical Encoding

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Soundness Problem (2): Unsound Case

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def bar(a: A):  
    # ensures result1 == result2  
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```

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def bar(a: A):  
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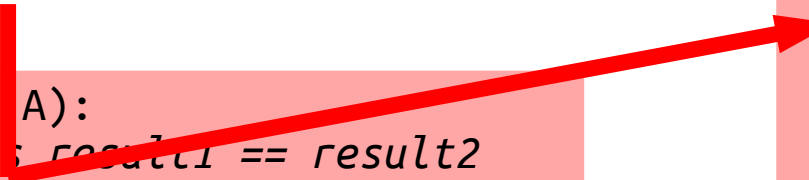
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UNSOUND

Criterion: Operationality

Encodings that encode *operational* behavior are always sound

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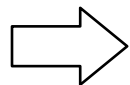
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$\langle C, \Sigma \rangle$

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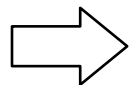


$\langle C', \Sigma' \rangle$

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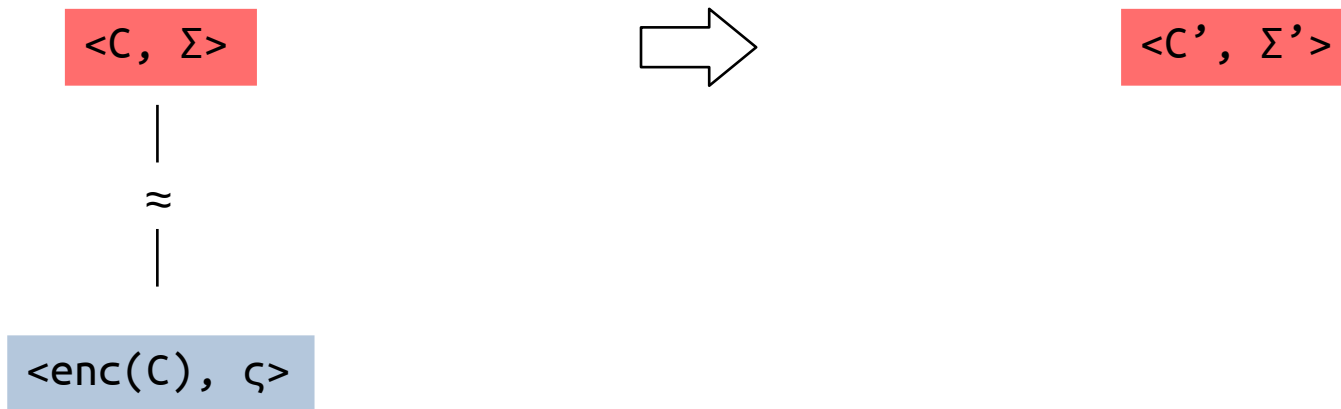


$\langle C', \Sigma' \rangle$

$\langle \text{enc}(C), \varsigma \rangle$

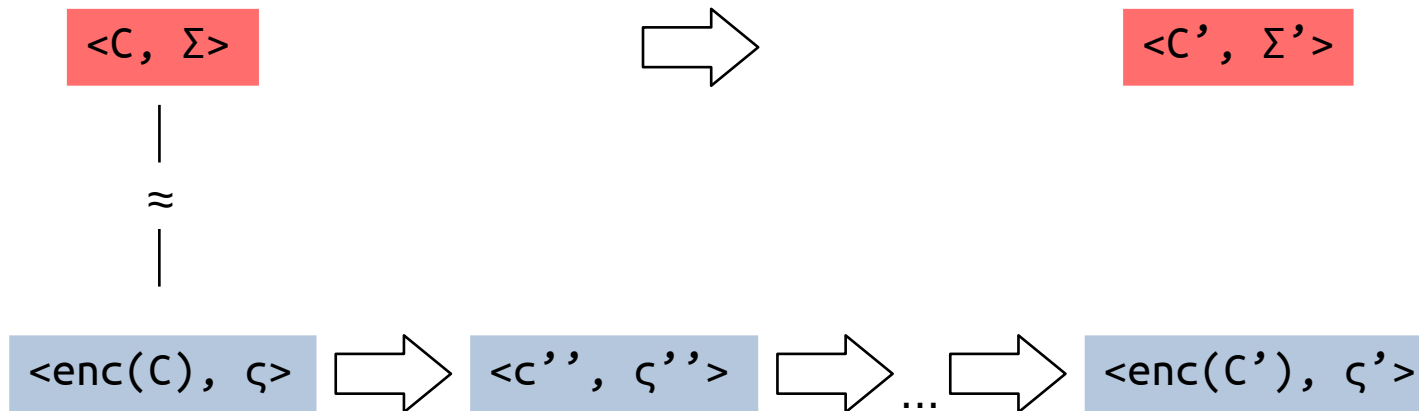
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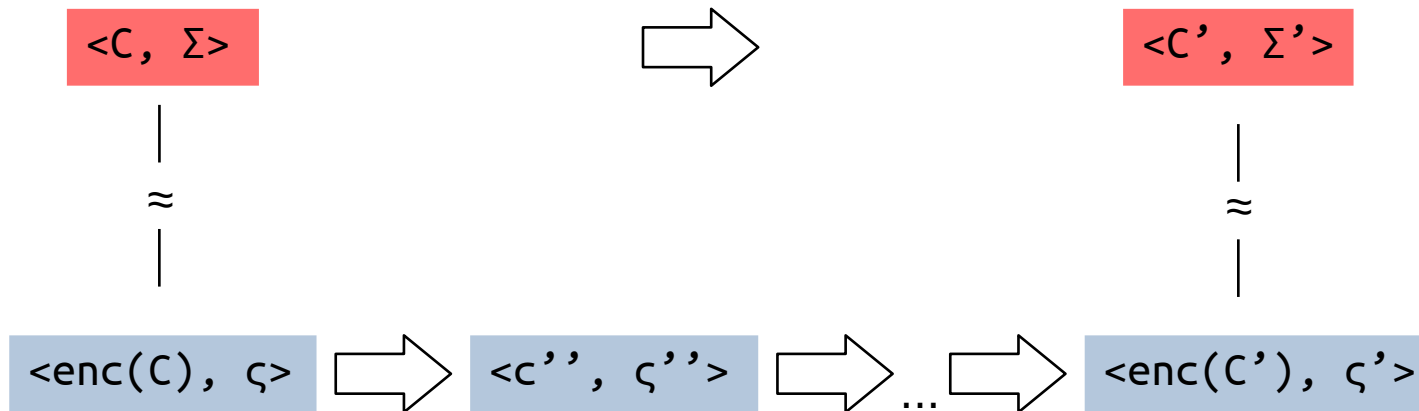
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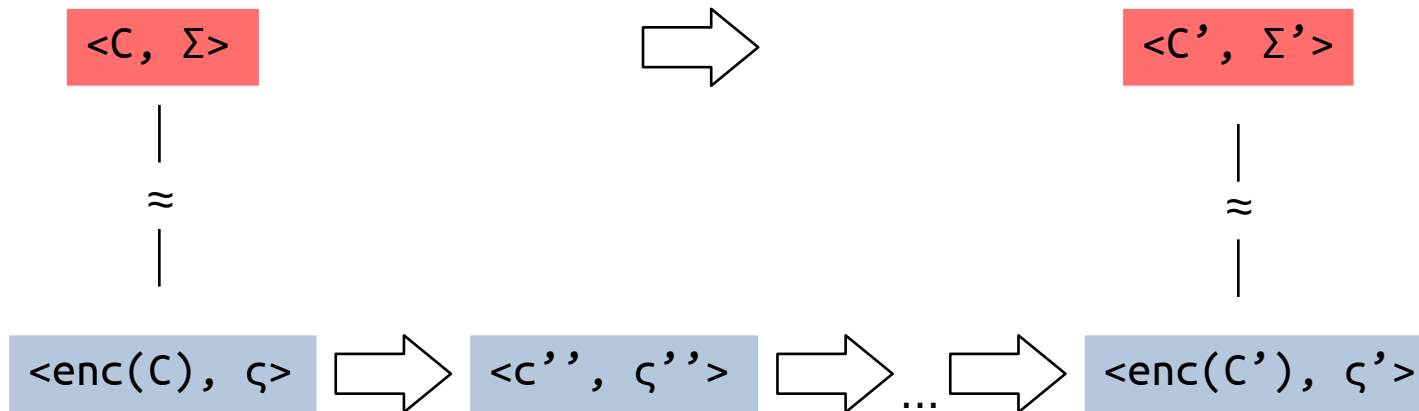
Encodings that encode *operational* behavior are always sound



Criterion: Operationality

Encodings that encode *operational* behavior are always sound

- Sufficient criterion, not necessary



Practical Relevance

- Typical: Frontend encodes *most* language features operationally
 - Main advantage of IVLs

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- Typical: Frontend encodes *most* language features operationally
 - Main advantage of IVLs
- *Easily identify* non-operational parts
 - Often intentional to ensure modularity
 - Must be checked case-by-case
- Examples of non-operational encodings:
 - Dynamically-bound calls (Dafny, Nagini)
 - Local blocks (VCC)
 - Loops (Prusti)

Soundness Problem (3): Possible Solution

```
def bar(a: A):  
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```

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assert  
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Soundness Problem (3): Possible Solution

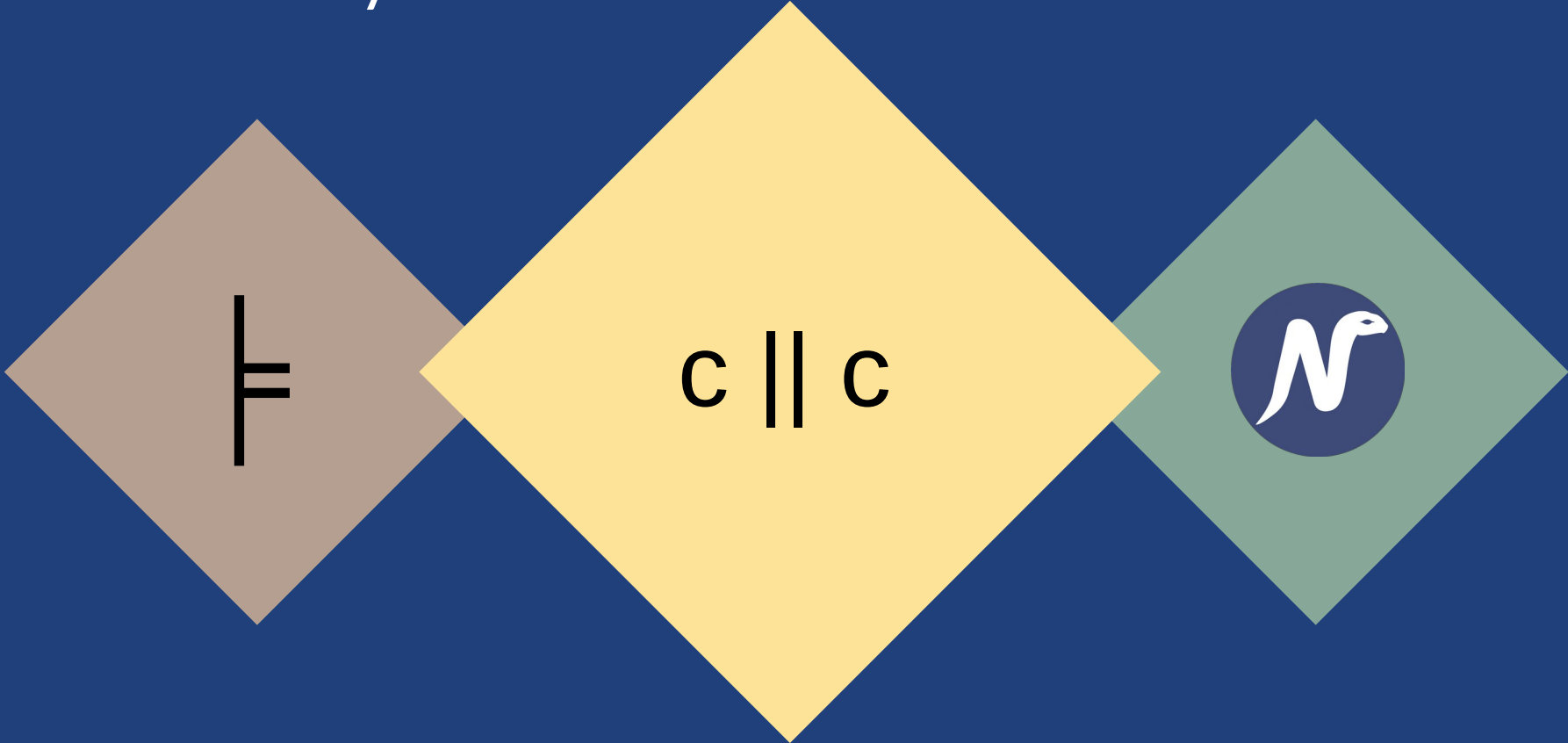
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def bar(a: A):  
    # ensures result1 == result2  
    return a.foo()
```

```
assert  
    (result1 == result2) ==>  
    (result1 == result2)
```

```
def bar(a: A):  
    # ensures result1 == result2  
    assume type(a1) == type(a2)  
    ==> result1 == result2
```

```
class A:  
    def foo(self):  
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Concurrency

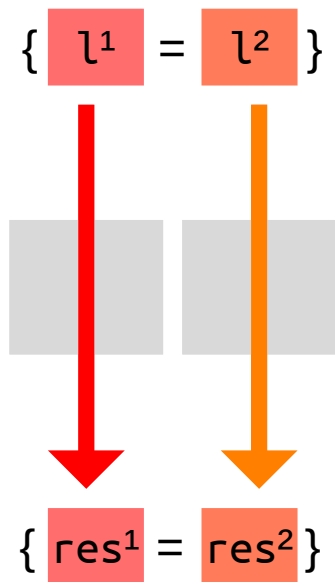


$\models$

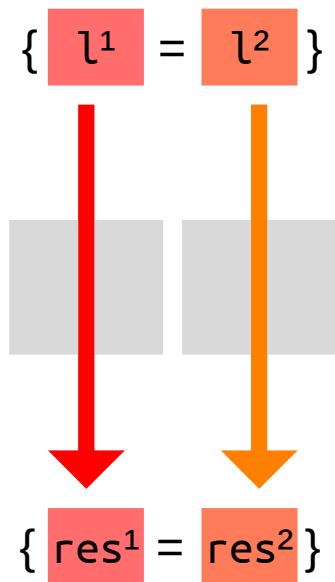
$c \parallel c$



Concurrency: Noninterference Properties

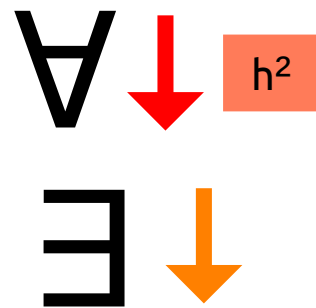
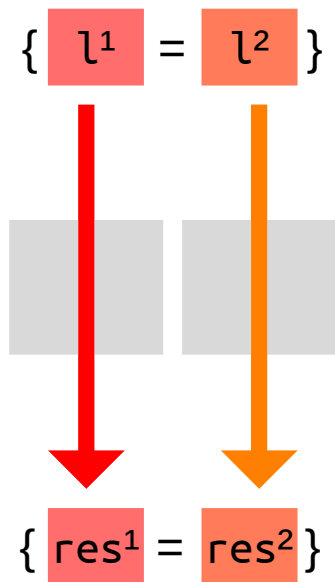


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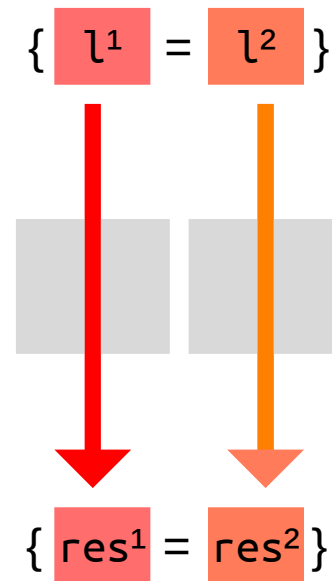
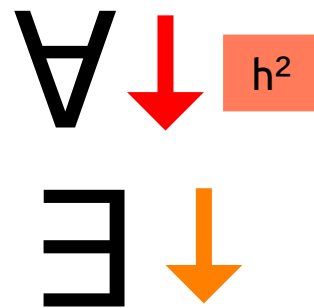
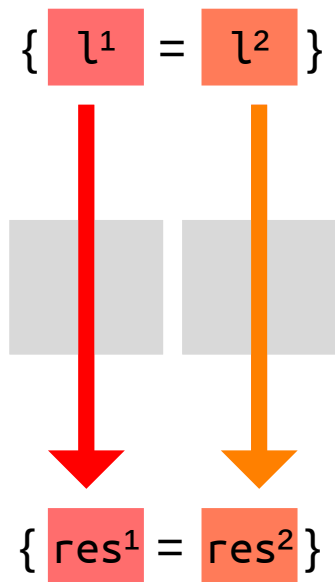
Noninterference: Too strict

Concurrency: Noninterference Properties



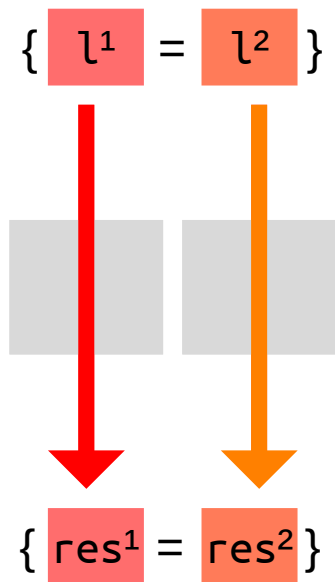
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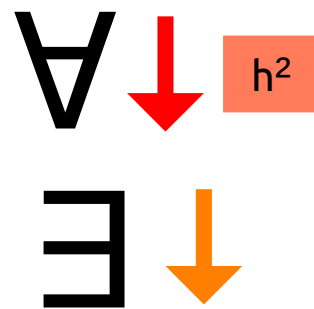


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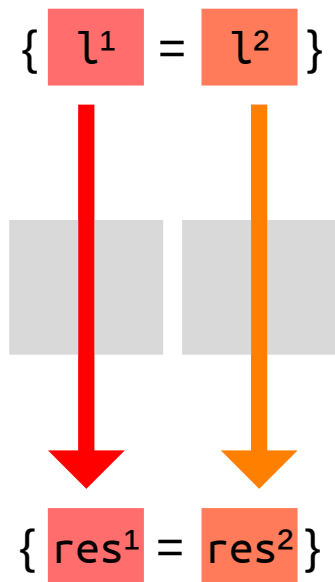


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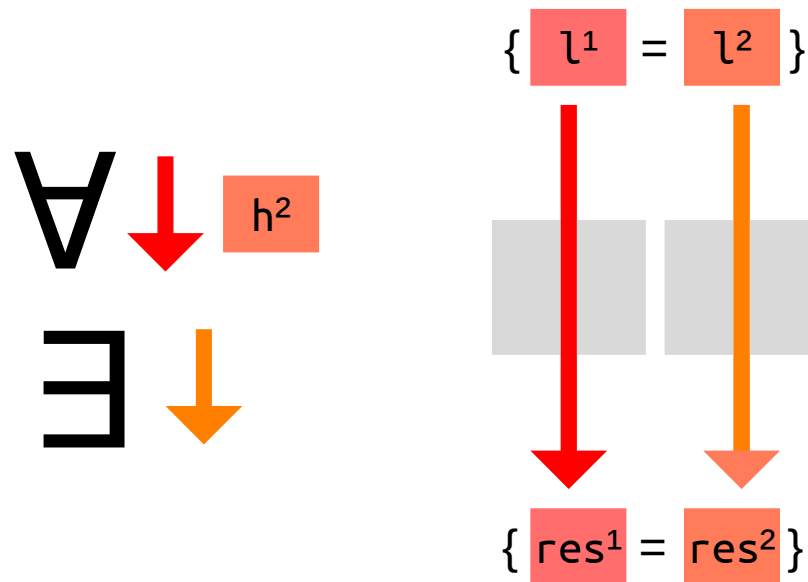


Possibilistic Noninterference

Concurrency: Noninterference Properties



Noninterference: Too strict



Possibilistic Noninterference
(Probabilistic Noninterference in the
paper)

Concurrency: Existing Encodings

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- Prove data race freedom

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 - Use concurrent separation logic / ownership / ...

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enc(s)

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- Encode checks into sequential IVL

Concurrency: Existing Encodings

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acquire l

enc(s)

<get ownership of l's heap memory>
assume Inv(l)

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Concurrency: Existing Encodings

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acquire l

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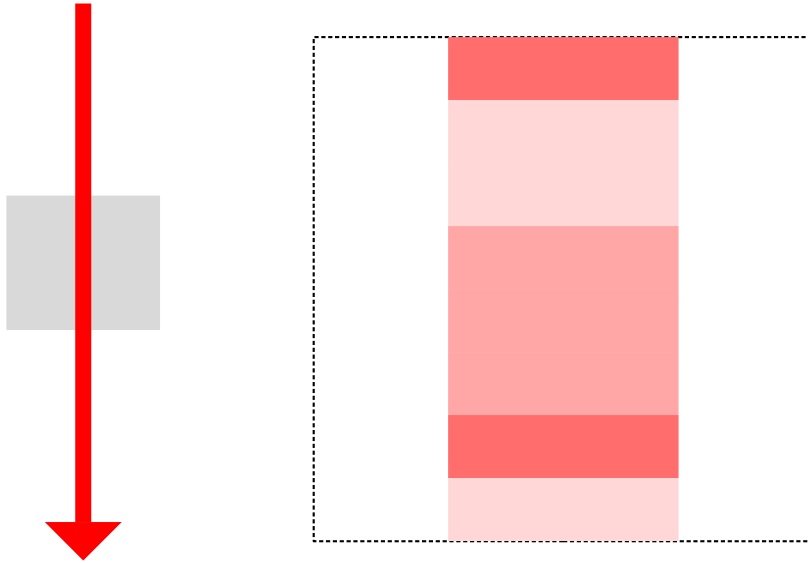
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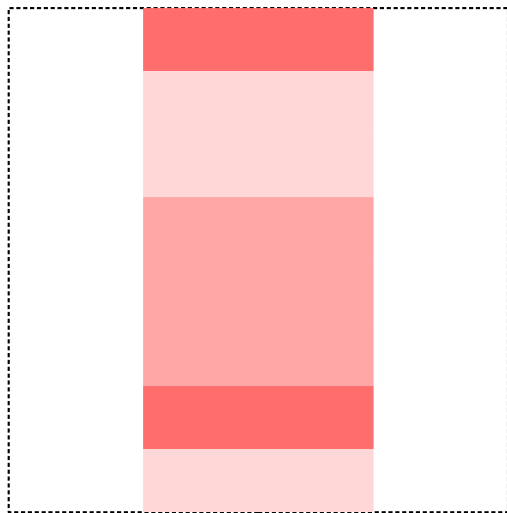
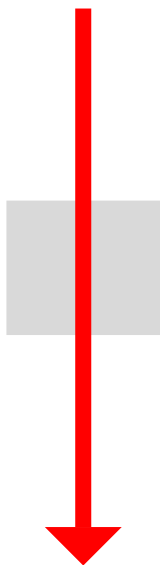
- Prove data race freedom
 - Use concurrent separation logic / ownership / ...
- Encode checks into sequential IVL
- Goal: adapt this encoding

Possibilistic Noninterference: Approach



Given first trace,

Possibilistic Noninterference: Approach

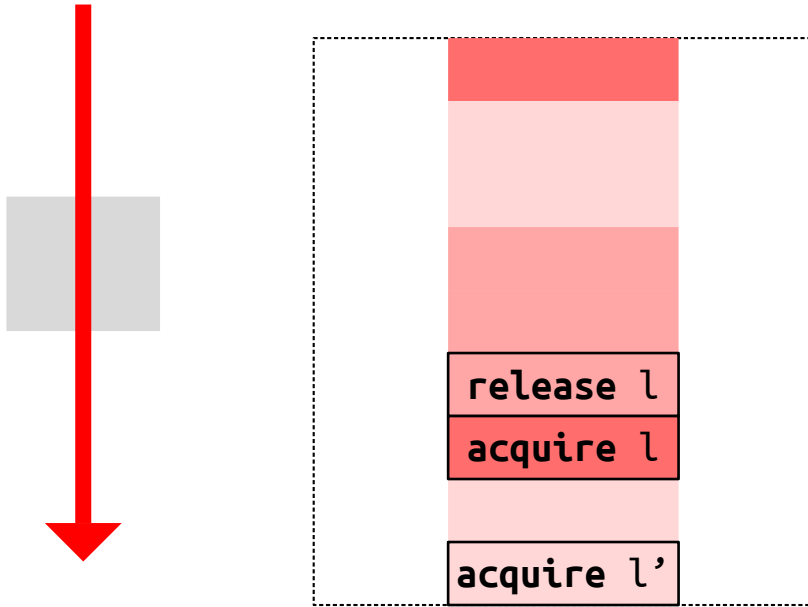


Given first trace,

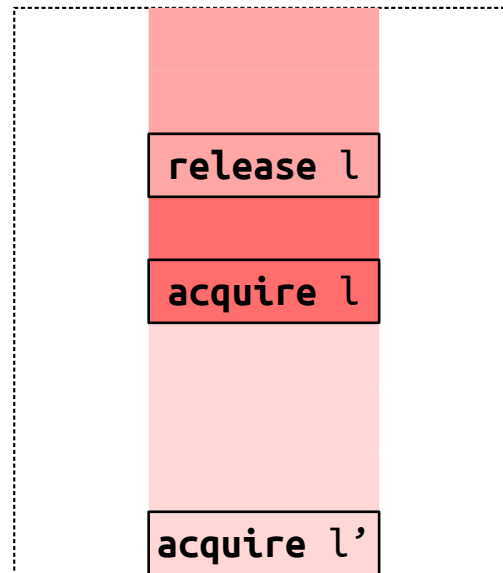
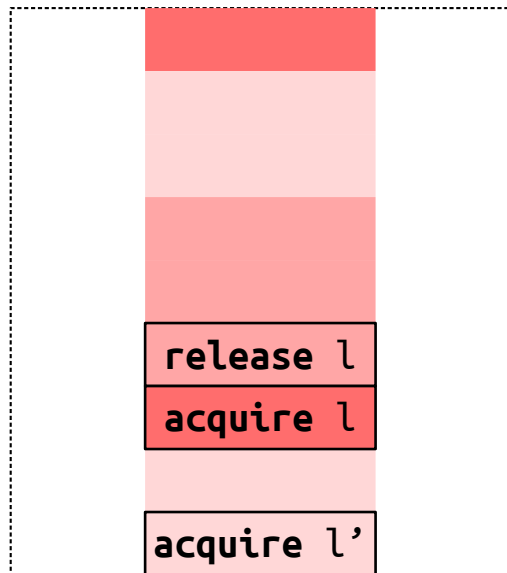
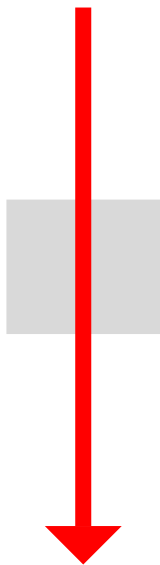


ensure second trace exists

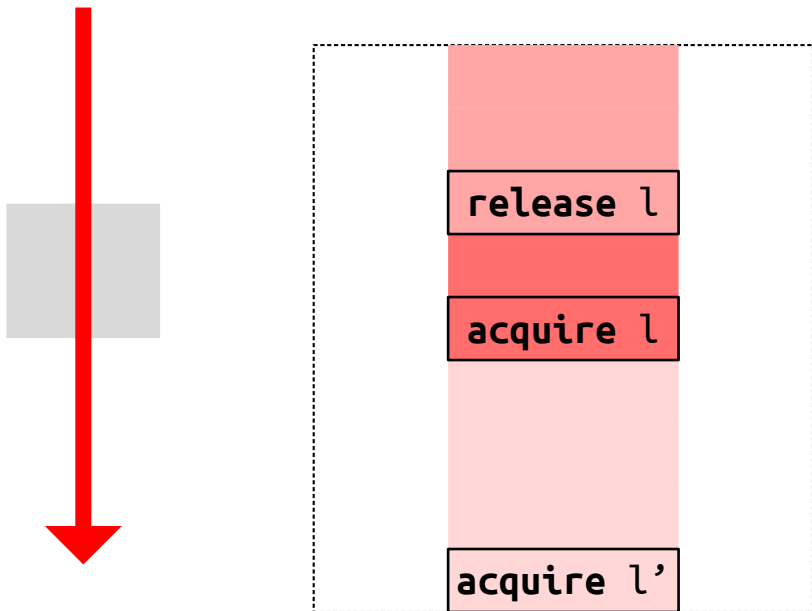
Possibilistic Noninterference: Simplification



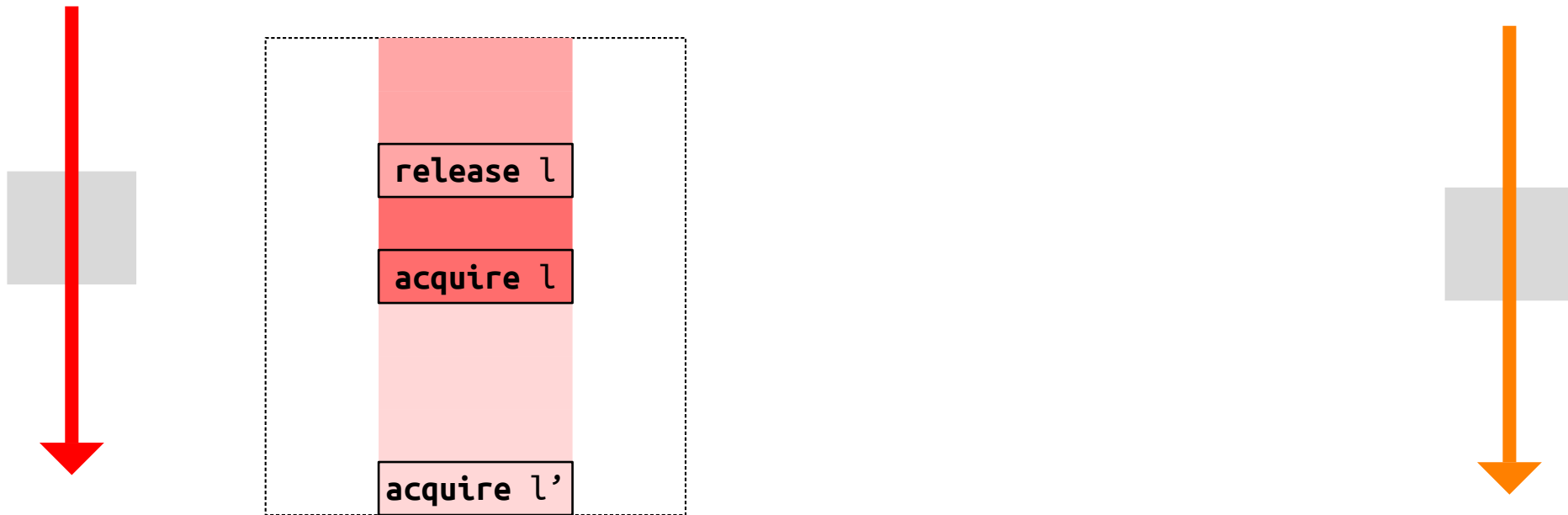
Possibilistic Noninterference: Simplification



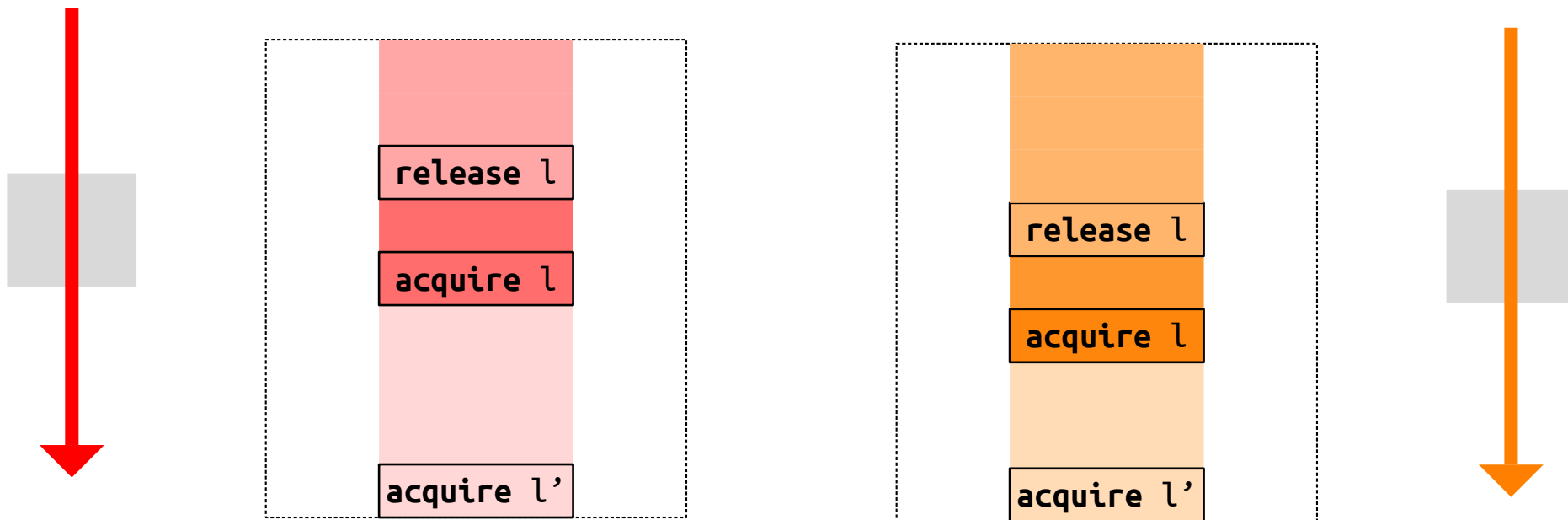
Possibilistic Noninterference: Goal



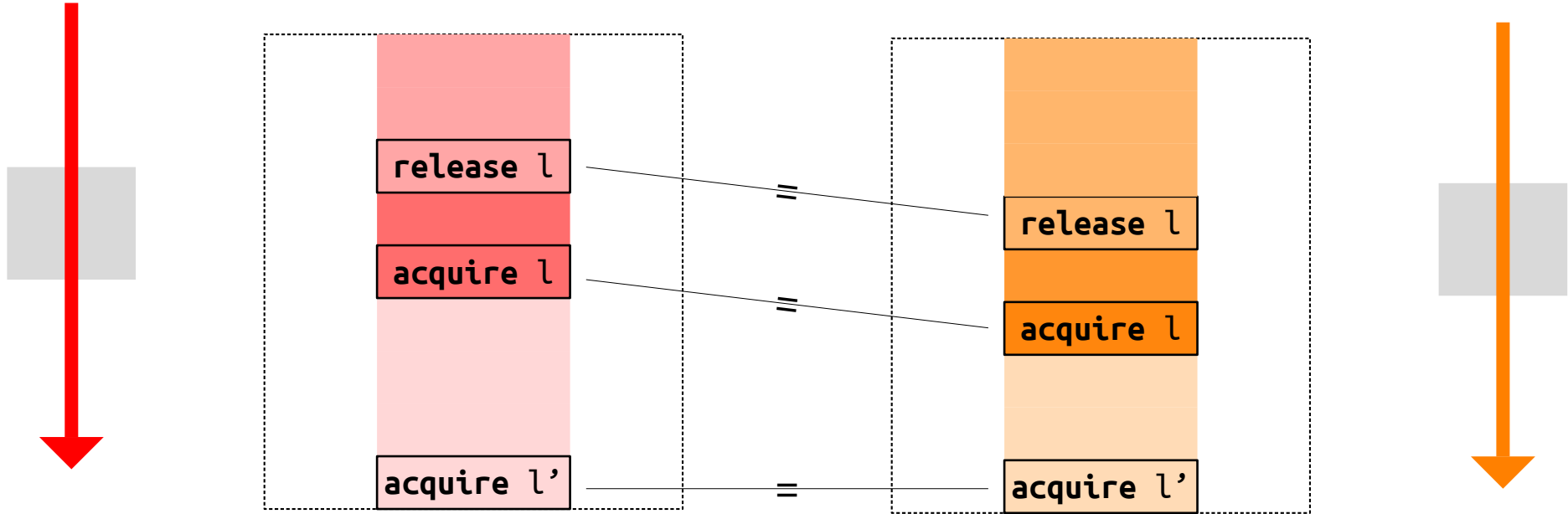
Possibilistic Noninterference: Goal



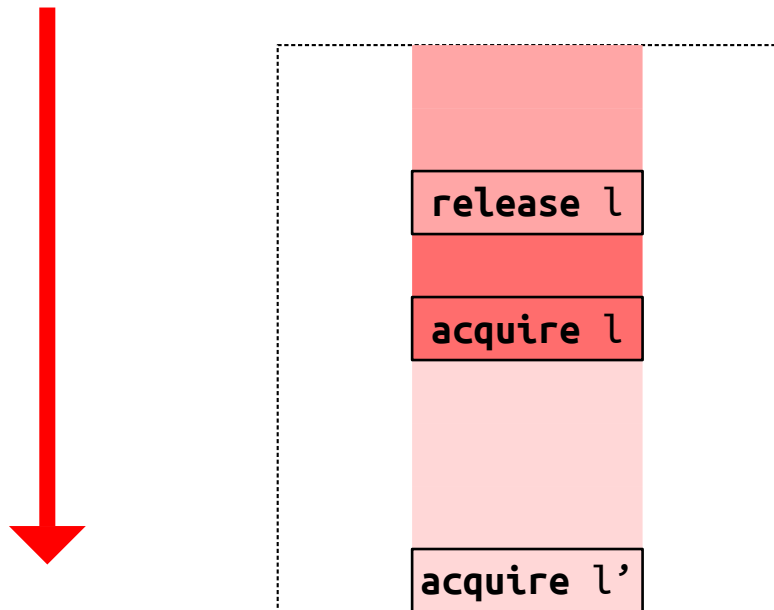
Possibilistic Noninterference: Goal



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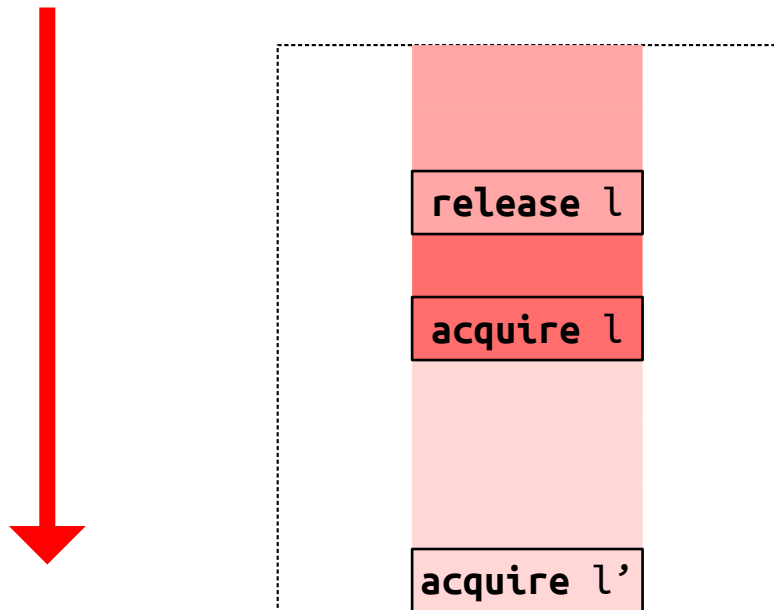


Possibilistic Noninterference: Conditions



Conditions for every lock operation:

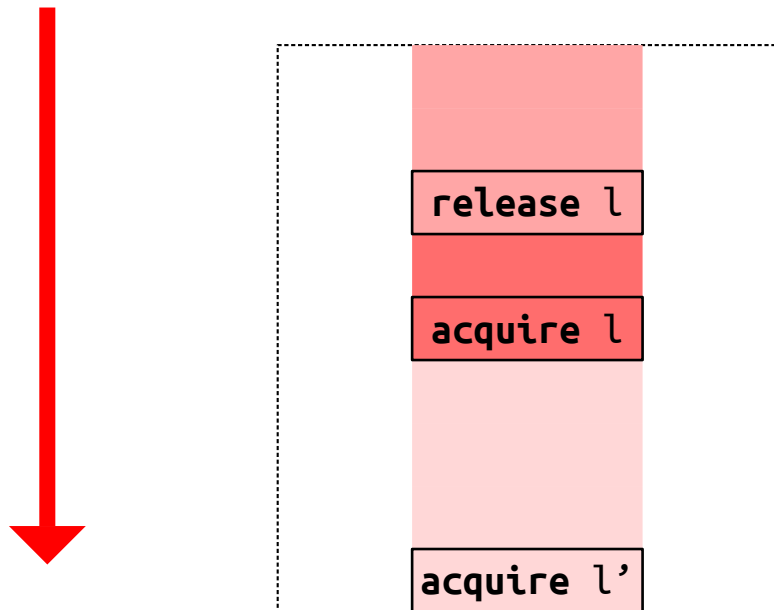
Possibilistic Noninterference: Conditions



Conditions for every lock operation:

- Not under high condition

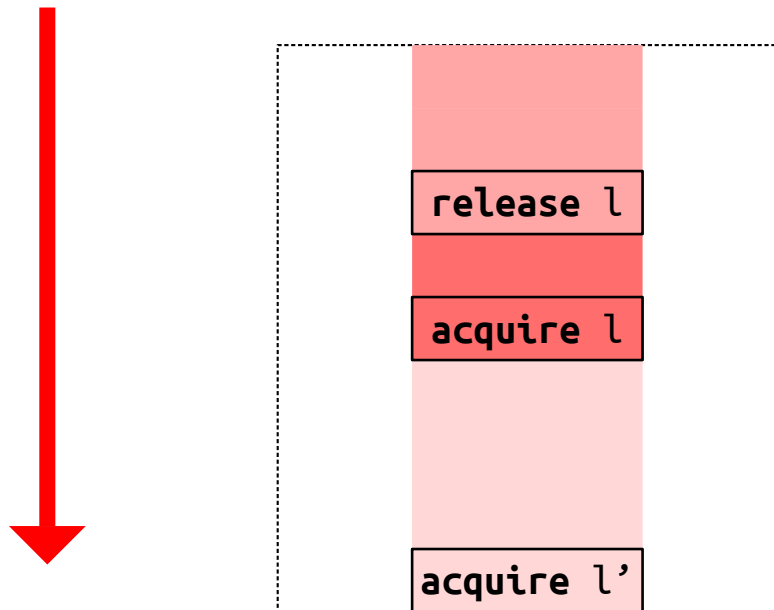
Possibilistic Noninterference: Conditions



Conditions for every lock operation:

- Not under high condition
- Code termination independent of high data

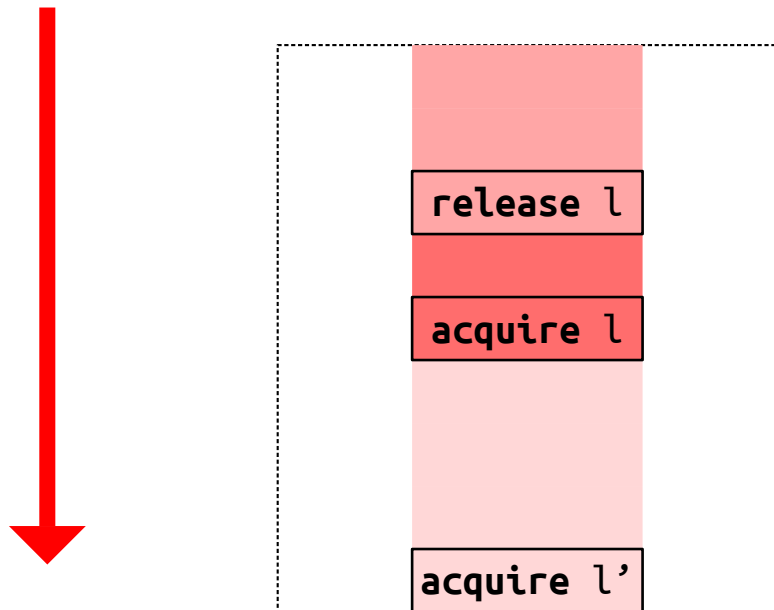
Possibilistic Noninterference: Conditions



Conditions for every lock operation:

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- Executed on same lock in both traces

Possibilistic Noninterference: Conditions



Conditions for every lock operation:

- Not under high condition
- Code termination independent of high data
- Executed on same lock in both traces
- Protected data identical in both traces

Possibilistic Noninterference: Encoding

- Not under high condition
- Code termination independent of high data
- Executed on same lock in both traces
- Protected data identical in both traces

s	enc'(s)
acquire l	<get ownership of l's heap memory> assume Inv(l)
release l	assert Inv(l) <give up ownership of l's heap memory>

Possibilistic Noninterference: Encoding

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- Code termination independent of high data
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Possibilistic Noninterference: Encoding

- Not under high condition
- Code termination independent of high data
- Executed on same lock in both traces
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s	enc'(s)
acquire l	assert <control flow is low> <get ownership of l's heap memory> assume Inv(l)
release l	assert <control flow is low> assert Inv(l) <give up ownership of l's heap memory>

Possibilistic Noninterference: Encoding

- Not under high condition
- Code termination independent of high data
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acquire l	assert <control flow is low> <get ownership of l's heap memory> assume Inv(l)
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while e do s	assert e1 == e2 enc(while e do s; assert e1 == e2)

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- Executed on same lock in both traces
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release l	assert <control flow is low> assert l1 == l2 assert Inv(l) <give up ownership of l's heap memory>
while e do s	assert e1 == e2 enc(while e do s; assert e1 == e2)

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Possibilistic Noninterference: Encoding

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- Code termination independent of high data
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s	enc'(s)
acquire l	assert <control flow is low> assert l1 == l2 <get ownership of l's heap memory> assume Inv(l1, l2)
release l	assert <control flow is low> assert l1 == l2 assert Inv(l1, l2) <give up ownership of l's heap memory>
while e do s	assert e1 == e2 enc'(while e do s; assert e1 == e2)

Possibilistic Noninterference: Example Violation

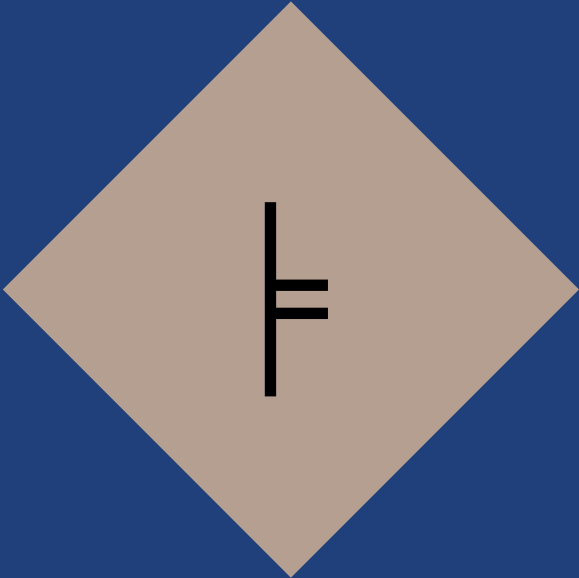
```
def main(h: bool):  
  c := new Cell()  
  l := new CellLock(c)  
  fork thread2(c, l)  
  c.val := 4  
  if h:  
    release l  
    acquire l  
  c.val := 5  
  release l
```


Possibilistic Noninterference: Example Violation

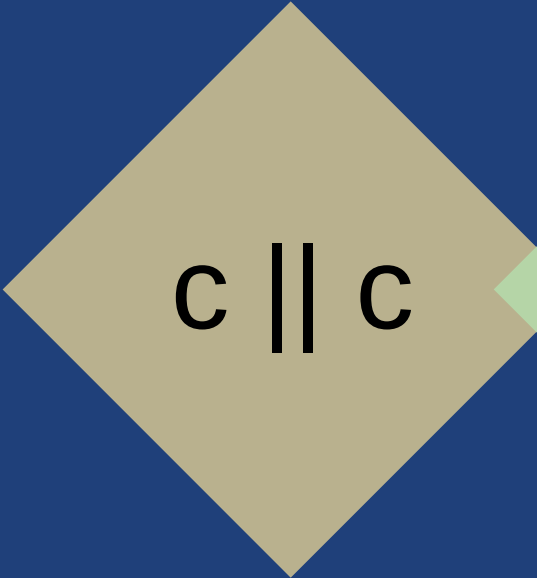
```
def main(h: bool):  
    c := new Cell()  
    l := new CellLock(c)  
    fork thread2(c, l)  
    c.val := 4  
    if h:  
        release l  
        acquire l  
    c.val := 5  
    release l
```

```
def thread2(c: Cell, l: CellLock):  
    acquire l  
    print(c.val)  
    # if 4 is printed, h is true  
    release l
```

Implementation and Evaluation



$\models$



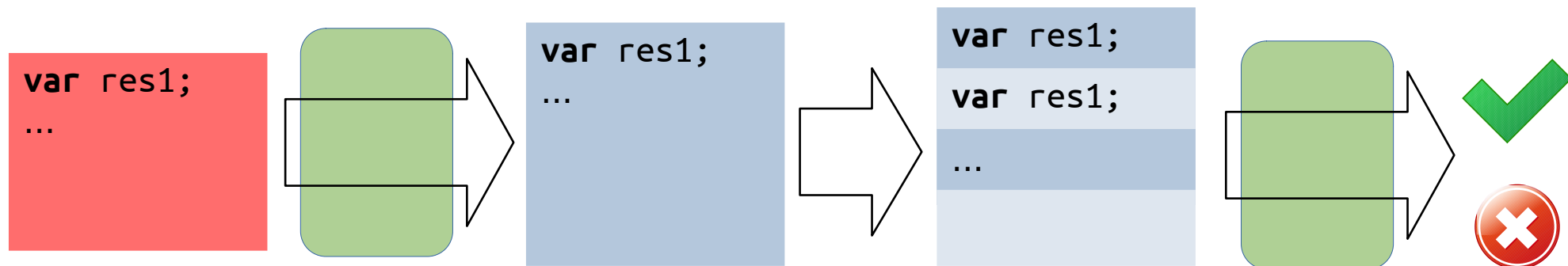
$c \parallel c$



Implementation in Nagini / Viper

[Eilers & Müller, CAV 2018]

[Müller, Schwerhoff, Summers, VMCAI 2016]



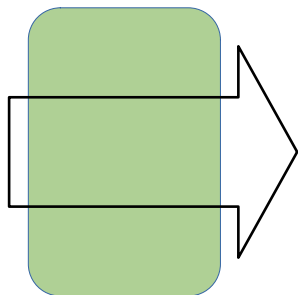
Implementation in Nagini / Viper

[Eilers & Müller, CAV 2018]

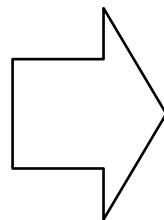
[Müller, Schwerhoff, Summers, VMCAI 2016]



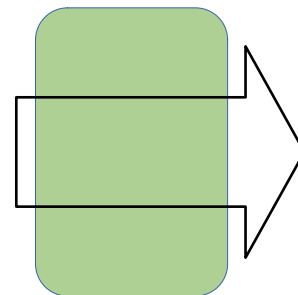
```
var res1;  
...
```



```
var res1;  
...
```



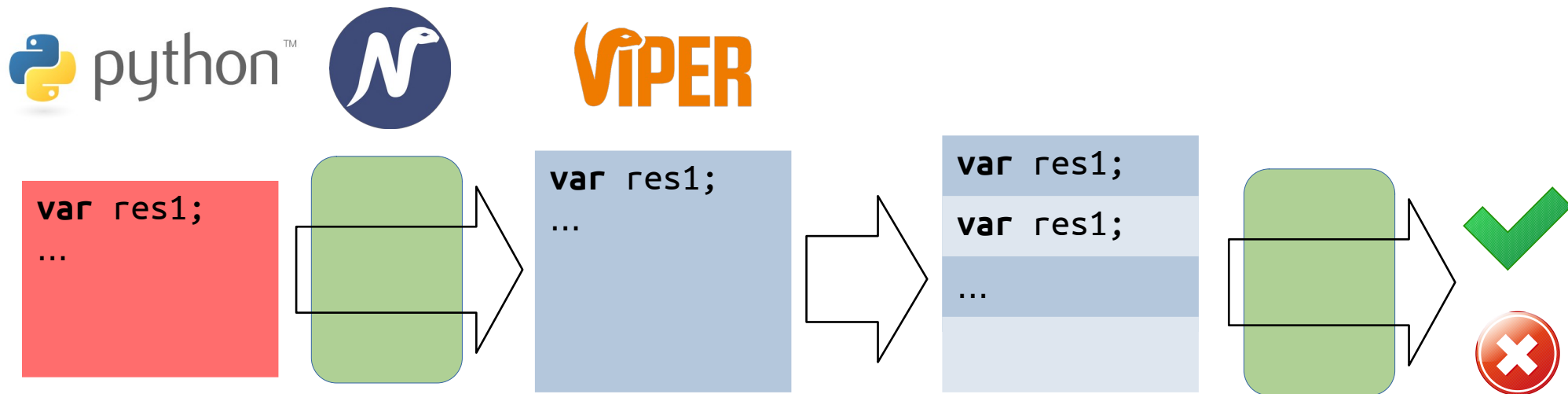
```
var res1;  
var res1;  
...
```



Implementation in Nagini / Viper

[Eilers & Müller, CAV 2018]

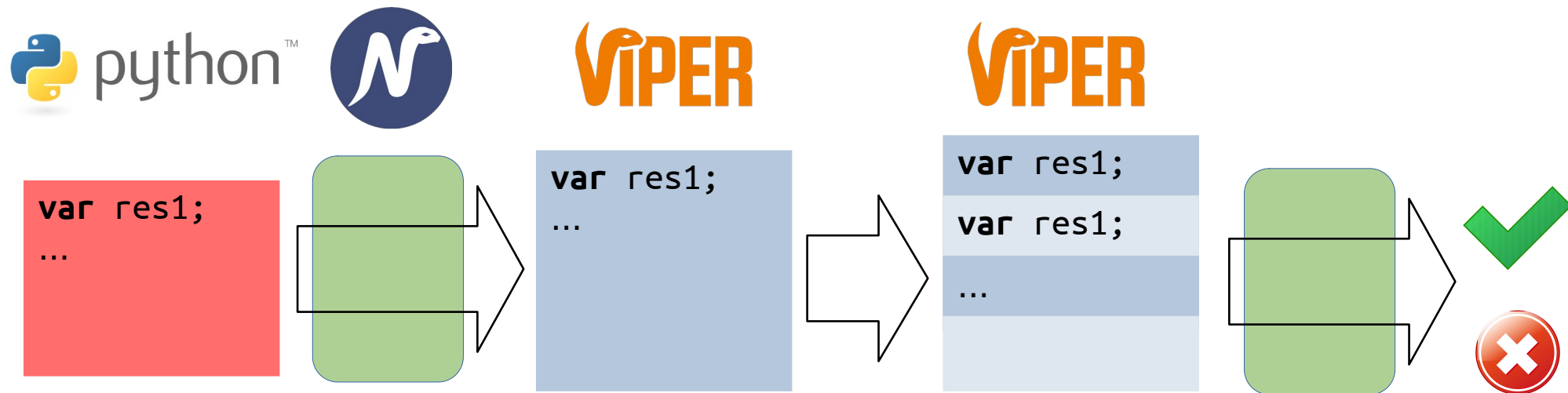
[Müller, Schwerhoff, Summers, VMCAI 2016]



Implementation in Nagini / Viper

[Eilers & Müller, CAV 2018]

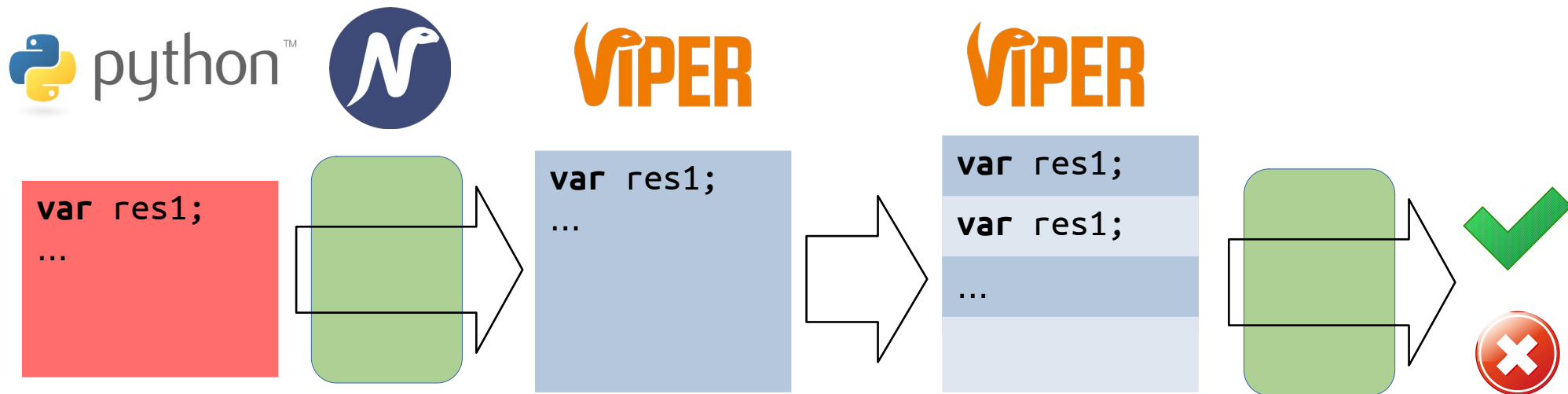
[Müller, Schwerhoff, Summers, VMCAI 2016]



Implementation in Nagini / Viper

[Eilers & Müller, CAV 2018]

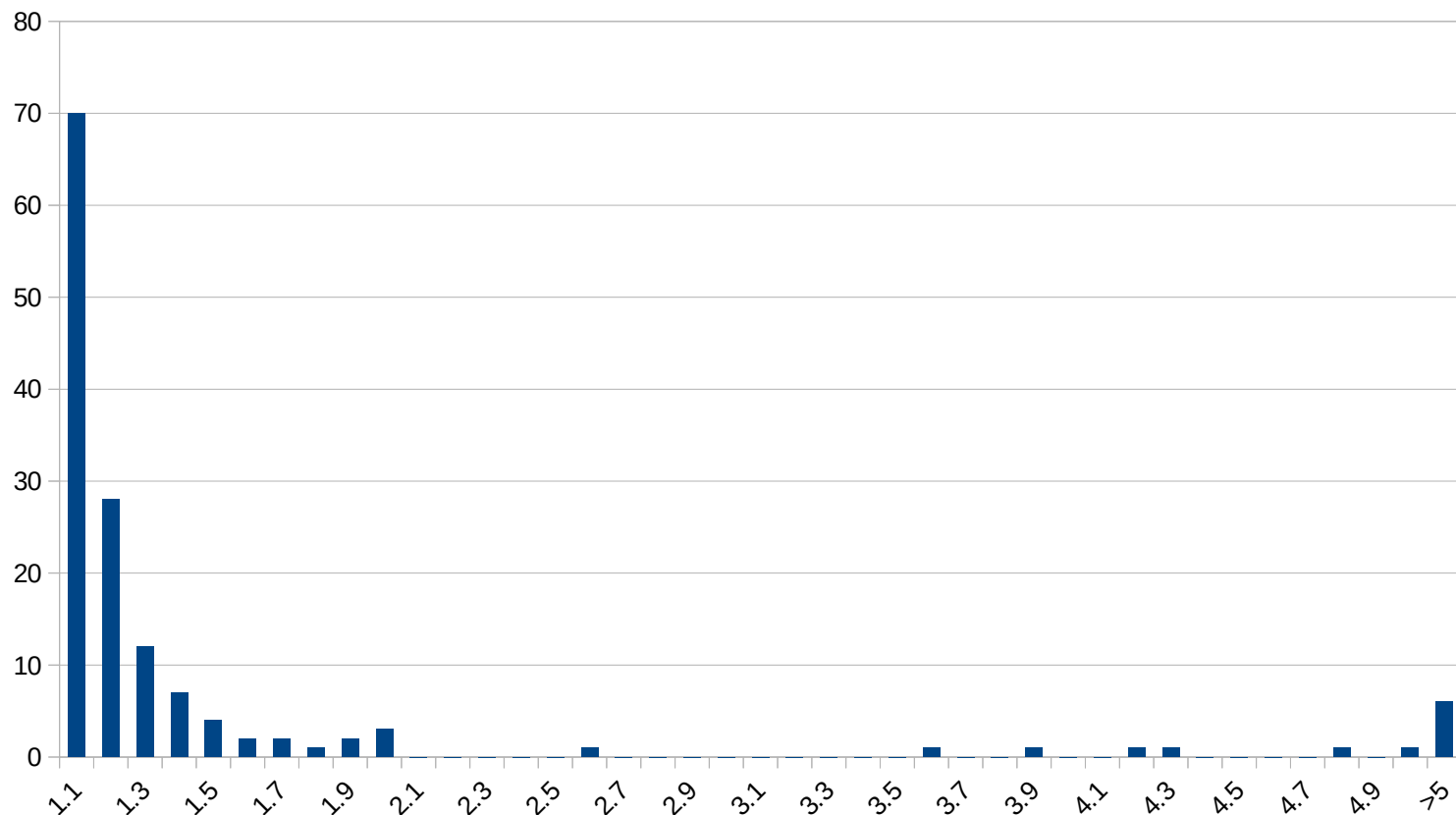
[Müller, Schwerhoff, Summers, VMCAI 2016]



Evaluation: Performance Overhead

- Slowdown introduced by product construction over test suite (144 tests)

Evaluation: Performance Overhead



- Slowdown introduced by product construction over test suite (144 tests)

Evaluation: Expressiveness

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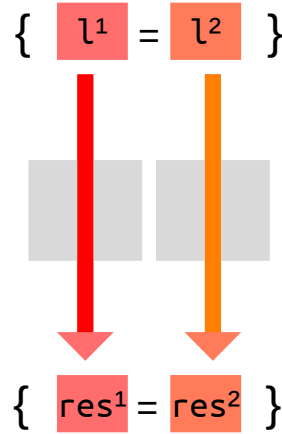
- 27 sequential examples from the literature
 - Declassification, termination-sensitive noninterference, value-dependent sensitivity

Evaluation: Expressiveness

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 - Declassification, termination-sensitive noninterference, value-dependent sensitivity
- Comparison with SecC [Ernst & Murray, CAV 2019]
- 5 examples from SecC test suite
 - Includes CDDC case study

Evaluation: Expressiveness

- 27 sequential examples from the literature
 - Declassification, termination-sensitive noninterference, value-dependent sensitivity
- Comparison with SecC [Ernst & Murray, CAV 2019]
- 5 examples from SecC test suite
 - Includes CDDC case study
- Nagini matches expressiveness of SecC
- SecC substantially faster
- Substantially lower development effort for extended Nagini version
 - Reused entire Python encoding and backend
 - Reusable product construction



assert

```
(result1 == result2) ==>
(result1 == result2)
```

def bar(a: A):

```
# ensures result1 == result2
assume result1 == result2
```

